





Navigation in a multi-obstacle environment. From partition of the space to a zonotopic-based MPC.

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# Outline

Preliminaries

Obstacle avoidance

Conclusions

# Outline





3 Conclusions

Navigation through multi-obstacle environments is one of the most challenging and intensively studied problem in the control and robotics communities.

The main difficulty: the non-convexity of the feasible regions in the motion space and consequently the lack of connectivity in the solution space.

State-of-the-art methods:

- optimization-based (Chen et al. 2016; Janeček et al. 2017; Szmuk et al. 2017)
- sampled(graph)-based (LaValle 2006; Weiss et al. 2017)



Main ideas:

- consider efficient mixed-integer descriptions of the non-convex region(s)
- generate feasible paths based on space partitioning

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## Generic control strategy

Model Predictive Control (MPC) (Mayne et al. 2000; Maciejowski 2002)

$$\mathcal{J}(x,u) = \left( \|x_{k+N_{P}|k} - \bar{x}_{\mathrm{ref}|k}\|_{P}^{2} + \sum_{l=1}^{N_{P}-1} \|x_{k+l|k} - \bar{x}_{\mathrm{ref}|k}\|_{Q}^{2} + \sum_{l=0}^{N_{P}-1} \|\Delta u_{k+l|k}\|_{R}^{2} \right)$$

Ingredients:

- quadratic/(non)linear optimization criterion
- (internal) model
- state and input constraints
  - magnitude constraints
  - obstacle avoidance constraints

 $\begin{aligned} x_{k+l|k} &= A x_{k+l-1|k} + B u_{k+l-1|k}, \\ x_{k+l|k} &\in \mathcal{X}, u_{k+l|k} \in \mathcal{U}, \\ x_{k+l|k} \notin \mathbb{P}. \end{aligned}$ 





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#### Preliminaries

#### Obstacle avoidance

- Zonotopes
- Space Partitioning
- Geometric path generation

#### Conclusions

## **Zonotopic representation**

Idea: Use zonotopes to characterize the regions of interest (Althoff et al. 2010; Stoican et al. 2013)

A zonotope is a centrally symmetric polytope and can be defined as a Minkowski sum of line segments (Kühn 1998)

$$Z = \{\mathbf{c} + \sum_{i=1}^{n_{\mathbf{g}}} \xi_i \mathbf{g}_i : \|\xi\|_{\infty} \le 1\}$$

 $Z = \mathcal{Z}(G, c) = \{G\xi + c | \xi \in \mathbb{R}^{n_g}, \|\xi\|_{\infty} \leq 1\}$ 

Several properties are of interest (Fukuda 2004).

-

- are closed under linear transformation and under Minkowski sum
- are symmetric,w.r.t. their centers
- their volume has an explicit formulation (Gover et al. 2010)

They are increasingly used in control applications due to their numerical robustness and simplicity, well-suited for large-scale problems (Althoff 2015).



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#### Zonotopes

## Zonotopic over-approximations I

In order to efficiently over-approximate the given shape of the obstacles, we can parametrize the zonotopes (Althoff et al. 2010), with respect to some fixed direction (an a priori given matrix  $G \in \mathbb{R}^{d \times m}$ ):

 $\mathcal{Z}(G\Delta_j, c_j), \quad j = 1 \dots N_o.$ 

where  $\Delta_{\ell}$  is a diagonal matrix. The k-th diagonal element is noted as  $\delta_{i_k}$ .

Consequences:

- the generated half-spaces share common normal vectors irrespective of the scaling factors.
- the resulting hyperplanes are parallel with each other
- simplified formulation



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## Zonotopic over-approximations II

Goal: Provide adequate zonotopic over-approximations for a multi-obstacle environment:



#₽	$\#\mathbb{H}$	$\gamma^*(N)$	$t_{\gamma(N)}(sec)$	$\#\Sigma_{\mathbb{P}}$
7	34	419	9.22	75

$$\begin{split} (\Delta_j, c_j)^* =& rg\min_{\Delta_j, c_j} \mathcal{C}(\Delta_j, c_j) \ & ext{ s.t. } P_j \subseteq \mathcal{Z}(G\Delta_j, c_j) \end{split}$$

Measures  $C(\Delta_j, c_j)$ :

i)  $Vol(\mathcal{Z}(G\Delta_j, c_j))$  -zonotope volume

ii) 
$$\|\delta\|_1$$
 -generator sum  $\left(\sum_{k=1}^m g_k \delta_{j_k}\right)$ 

iii) 
$$\left\|\delta\right\|_{\infty}$$
 -largest generator  $(\max_{k=1\dots m} g_k \delta_{j_k})$ 

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$$\begin{split} (\Delta_j, c_j)^* &= \arg\min_{\Delta_j, c_j} \mathcal{C}(\Delta_j, c_j) \\ &\text{s.t.} \ \left| h_i^\top (v_k - c_j) \right| \leq k_i(\Delta_j), \ \forall k. \end{split}$$

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## Illustrative example for zonotopic representation I





Figure: Volume with  $G = G_3$ .

$$G_3 = \begin{bmatrix} -1 & 0 & -1 \\ 0 & 1 & 1 \end{bmatrix}$$

Figure: Volume with  $G = G_4$ .

$$G_4 = egin{bmatrix} 1 & 0 & 0 \ 0 & 1 & 0 \ 0 & 0 & 1 \end{bmatrix}$$

## Illustrative example for zonotopic representation II

	Measure	G	t <sub>sol</sub>	$\#\mathbb{H}$	$\gamma^*(N)$	$\frac{\Delta\gamma(N)}{\gamma(N)}$	$t_{\gamma^*(N)}(\sec)$	$\#\Sigma_{\mathbb{P}}$	V	$\frac{\Delta V}{V}(\%)$
	$\ \delta\ _1$	<b>G</b> 1	8,13	42	505	20,53	9,53	197	376,98	71,7
		<b>G</b> <sub>2</sub>	8,02	28	225	-46,30	3,81	101	410,07	86,78
		G <sub>3</sub>	8,27	42	534	27,45	10,09	167	368,93	68,04
	$\ \delta\ _{\infty}$	<i>G</i> <sub>1</sub>	8,19	42	441	5,25	8,19	374	897,92	308,98
d = 2		G <sub>2</sub>	8,01	28	225	-46,30	3,91	175	583,33	165,69
		<i>G</i> <sub>3</sub>	8,19	42	441	5,25	8,19	374	897,92	308,98
	Vol	<i>G</i> <sub>1</sub>	9,40	42	510	21,72	9,66	199	368,41	67,8
		<b>G</b> <sub>2</sub>	9,19	28	225	-46,30	3,75	101	410,07	86,78
		<i>G</i> <sub>3</sub>	9,30	40	530	26,49	10,27	169	374,99	70,8
d = 3	$\left\ \delta\right\ _{1}$	G4	9,82	60	8400	-85,31	105,98	934	1857,46	323,56
		G <sub>5</sub>	10,50	120	62480	9,26	1145,09	4952	2019,34	360,47
	$\ \delta\ _{\infty}$	G <sub>4</sub>	9,71	60	8000	-86,01	100,22	1127	2623,3	498,19
		G <sub>5</sub>	10,59	120	51396	-10,12	932,12	9432	5852,55	1234,56
	Vol	G <sub>4</sub>	11,02	60	8400	-85,31	105,07	934	1857,46	323,56
		G <sub>5</sub>	11,79	84	24528	-57,11	413,07	2218	1908,2	335,13

The fixed directions/Generators:

$$G_{1,2,3} \in \left\{ \begin{bmatrix} 1 & 0 & 1 \\ 0 & 1 & 1 \end{bmatrix}, \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}, \begin{bmatrix} -1 & 0 & -1 \\ 0 & 1 & 1 \end{bmatrix} \right\}, \ G_{4,5} \in \left\{ \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}, \begin{bmatrix} 1 & 0 & 0 & 1 \\ 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 1 \end{bmatrix} \right\}.$$

## Obstacle avoidance: polytopic vs. zonotopic representation





Figure: Polytopes P.

Figure: Zonotopes  $\mathbb{Z}$ .

	Topology	N <sub>goal</sub>	$t_{goal}(sec)$	$t_{worst}$ (sec)
4 - 2	₽	143	11.64	0.22
u – 2	Z	146	10.07	0.18
d _ 2	P	98	83.87	0.81
<i>u</i> = 5	Z	132	57.07	0.42

Table: MPC parameters:  $N_p = 10$ ,  $P = 10I_{2d}$ ,  $Q = I_{2d}$ ,  $R = I_d$ .





"How can we approximate the obstacles with zonotopic sets while, simultaneously, safeguarding the feasible paths of the initial problem?"

# Zonotopic approximations with corridors

Proposed solution: The inclusion of a separation hyperplane (Boyd et al. 2004) in the optimization problem:  $\mathcal{H}_{sep} = \{x \in \mathbb{R}^d : h_{sep}^\top x = k_{sep}\}$ , with its corresponding half-spaces  $\mathcal{R}_{sep}^+$  and  $\mathcal{R}_{sep}^-$ .

- adding a linear constraint:  $\mathcal{Z}(G\Delta, c) \subset \mathcal{R}_{sep}^{\pm}$
- Shortcoming: usually, infeasible

- adding generators spanning  $\mathcal{H}_{sep}$
- Shortcoming: no inclusion monotonicity



#### Theorem

In  $\mathbb{R}^d$ , the maximum number of joint constraints for corridors with feasibility guarantees is d + 1.

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## **Space Partitioning**

Idea: Provide a partitioning of the navigation space w.r.t. the obstacles  $\mathbb{P} = \bigcup_{i=1}^{N_o} P_i$ 



#### Definition

A family of sets  $\{X_i\}_{i \in \mathcal{I}}$  verifying:

i) 
$$\mathbb{X} = \bigcup_{i=1}^{N_o} X_i$$
,

ii) 
$$\operatorname{int}(X_i) \cap \operatorname{int}(X_j) = \emptyset, \forall i \neq j \in \mathcal{I},$$

iii)  $P_i \subset int(X_i), \forall i \in \{1..., N_o\}$ 

is called a partition of  $\mathbb X$  induced by the obstacles  $\mathbb P.$ 

#### An intuitive solution

Idea: As  $P_i \cap P_j = \emptyset$ ,  $\forall i \neq j$ , the separating hyperplanes(Boyd et al. 2004) are candidates for the supporting hyperplanes of  $X_i \supset P_i$ 



 $<sup>\</sup>mathbb{X} \neq \bigcup_{i=1}^{N_o} X_i$ 

#### Alternatives:

- Generalized Voronoi Diagram (Afonso et al. 2013) or (Sugihara 1993)
- grid of square/cubic cells (Wang et al. 2015)
- convex lifting (Nguyen et al. 2018)

<sup>&</sup>lt;sup>1</sup>Nguyen, N. A., Gulan, M., Olaru, S., Rodriguez-Ayerbe, P. (2018). Convex lifting: Theory and control applications. IEEE Transactions on Automatic Control, 63(5), 1243-1258.

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## **Convex lifting-overview**

Definition (Convex lifting Nguyen et al. 2018)

For a polyhedral partition  $\{X_i\}_{i \in \mathcal{I}}$  of a domain  $\mathcal{X}$ , a piecewise affine lifting described by the function:

 $z(x) = a_i^\top x + b_i, x \in X_i,$ 

is called a convex (piecewise affine) lifting, if z(x) is continuous and convex over partition  $\{X_i\}_{i \in \mathcal{I}}$  of  $\mathcal{X}$ .

$$\begin{split} \min_{\mathbf{a}_i, b_i} \quad & \sum_{i=1}^{N_o} J(\mathbf{a}_i, b_i) = \sum_{i=1}^{N_o} \begin{bmatrix} \mathbf{a}_i^\top & b_i \end{bmatrix} \begin{bmatrix} \mathbf{a}_i \\ b_i \end{bmatrix} \\ \text{s.t.} \quad & \mathbf{a}_i^\top \mathbf{v} + b_i \leq M, \forall \mathbf{v} \in \mathcal{V}(X_i), \forall i, \\ & \mathbf{a}_j^\top \mathbf{v} + b_j \geq \mathbf{a}_i^\top \mathbf{v} + b_i + \epsilon, \forall \mathbf{v} \in \mathcal{V}(X_j) \setminus \mathcal{V}(X_i), \forall i \neq j, \\ & \mathbf{a}_i^\top \mathbf{v} + b_i = \mathbf{a}_j^\top \mathbf{v} + b_j, \forall \mathbf{v} \in \mathcal{V}(X_i \cap X_j), \forall i \neq j. \end{split}$$



$$\mathcal{P} = \left\{ \begin{bmatrix} x \\ z \end{bmatrix} \in \mathbb{R}^{d+1} : \begin{bmatrix} a_i^\top - 1 \end{bmatrix} \begin{bmatrix} x \\ z \end{bmatrix} \le -b_i, \ i \in \mathcal{I} \right\}.$$

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$$\begin{split} \min_{a_i,b_i} \quad & \sum_{i=1}^{N_o} J(a_i,b_i) = \sum_{i=1}^{N_o} \begin{bmatrix} a_i^\top & b_i \end{bmatrix} \begin{bmatrix} a_i \\ b_i \end{bmatrix} \\ \text{s.t.} \quad & a_i^\top v + b_i \leq M, \forall v \in \mathcal{V}(P_i), \forall i, \\ & a_i^\top v + b_i \geq a_j^\top v + b_j + \epsilon, \forall v \in \mathcal{V}(P_i), \forall i \neq j, \end{split}$$



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## From Lifting to Partitioning







## Space partitioning



## Geometric path generation

Idea: Define a graph  $\Gamma = (\mathcal{N}, \mathcal{E}, f), f : \mathcal{E} \to \mathbb{R}$ , based on the partition  $\{X_i\}_{i=1:N_0}$ .

- $\mathcal{N} = \{\mathcal{V}(X_i)\}_{i=1:N_o}$
- $\mathcal{E}$  the facets of the partition cells.



 $N_o = 31, d = 2$ 

 $N_o = 5, d = 3$ 

#### Next steps:

- connect  $x_0, x_f \in \mathcal{C}_{\mathbb{X}}(\mathbb{P})$  to  $\Gamma$
- run a graph search algorithm, e.g. Dijkstra's Algorithm (Karaman et al. 2011)

## Monte Carlo study for feasible path generation

Computing times w.r.t. the number of obstacles



### Illustrative example for obstacle avoidance

- $\{X_i\}_{i=1:N_o}$  and  $\Gamma$
- "local" zonotopic over-approximation

- Path $(x_i, x_f)$
- MPC with  $N_p = 7$

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#### Conclusions:

- we studied the benefits and difficulties of choosing a particular family of sets (parametrized zonotopes) for the non-convex feasible domain representation
- · we introduced a partitioning procedure of the workspace based on convex lifting
- we propose a navigation strategy with obstacle avoidance guarantees using local zonotopic approximations of the obstacles

#### Future directions:

- extension of the preliminary results for more complex scenarios (e.g., mobile obstacles)
- improvements of the MPC problem (feasibility etc.)
- develop a complete navigation strategy based on space partitioning

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#### Zonotopic sets – properties

Several properties are of interest (Fukuda 2004).

Let  $Z_1 = \mathcal{Z}(\mathbf{G}_1, \mathbf{c}_1) \subset \mathbb{R}^n$ ,  $Z_1 = \mathcal{Z}(\mathbf{G}_2, \mathbf{c}_2) \subset \mathbb{R}^n$ :

- are closed under linear transformation:  $R\mathcal{Z}(G_1, c_1) = \mathcal{Z}(RG_1, Rc_1);$
- are closed under Minkowski sum:  $\mathcal{Z}(\mathbf{G}_1, \mathbf{c}_1) \oplus \mathcal{Z}(\mathbf{G}_2, \mathbf{c}_2) = \mathcal{Z}(\begin{bmatrix} \mathbf{G}_1 & \mathbf{G}_2 \end{bmatrix}, \mathbf{c}_1 + \mathbf{c}_2);$
- $\bullet\,$  are symmetric, w.r.t. their center:  $-Z_1=-\mathcal{Z}\{\textbf{G}_1,\textbf{c}_1\}=\mathcal{Z}\{\textbf{G}_1,-\textbf{c}_1\}$
- their volume has an explicit formulation (Gover et al. 2010):  $\operatorname{Vol}(\mathcal{Z}(G, c)) = \sum_{1 \leq j_1 < j_2 \dots , j_d \leq m} \left| \det(G^{j_1 \dots j_d}) \right|$
- their corresponding half-space representation (Althoff et al. 2010):

$$\mathcal{Z}(G,c) = \bigcap_{1 \leq j_1 < \dots + j_{d-1} \leq m} \{x \in \mathbb{R}^d : |h_i(x-c)| \leq k_i\},\$$

$$h_i \perp g_{j_l}, \forall j_l \in \{j_1 \dots j_{d-1}\}, \quad k_i = \sum_{j_l \notin \{j_1 \dots j_{d-1}\}} \left| h_i^\top g_{j_l} \right|$$

Back to Zonotopic Approximation